A New Method for Controlling and Maintaining Topology in Wireless Sensor Networks

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ABSTRACT

Topology Control is an essential technique in a wireless sensor network to extend the operational time of the sensor nodes. The goal of this technique is to maintain network connectivity and optimize performance metrics such as network lifetime and throughput. In this paper we presented a new method for controlling and maintaining topology in wireless sensor networks that show some improvement over the state of art methods. The results are analyzed based on objective criteria.

KEYWORDS


1. INTRODUCTION

A wireless sensor network is a reliant collection of nodes that communicate over relatively bandwidth constrained wireless links. In general, topology control in sensor networks adjusts their transmission ranges to conserve energy and reduce interference. In case of wireless sensor networks the aspects which need to be considered while controlling topology are operation in hostile environments, Data processing and scalability. A significant amount of research works [1-4] have been done using non homogenous approaches such as location based, neighbor based and direction based.

In this paper we devised a new method called ENR which stands for Energy-Neighbour-Range to control and maintain topology in wireless sensor networks. Our approach is fully distributed, asynchronous and generates a connected topology. The method is based on the principle of maintaining the number of physical neighbours, transmission range and energy levels at each node while constructing topology in order to improve the efficiency of the network.

The rest of the paper is organized as follows. Section 2 briefly describes the taxonomy of topology control methods. Section 3 presents our working principle. Section 4 presents the simulation environment and experimental results. Finally conclusions are drawn in Section 5.

2. TAXONOMY OF TOPOLOGY CONTROL

Topology control is an essential technique in Wireless sensor network which leads to reorganizing and managing of certain node parameters and modes of operation from time to time to modify the topology of the network with the goal of extending its lifetime while preserving important characteristics such as connectivity and coverage. Topology control is an iterative process which consists of two phases those are topology construction and topology maintenance.
During the topology construction phase, each node builds a new topology based on the best parent node. Nodes can extend the network lifetime by adjusting transmission power to control the topology. In the topology maintenance phase, each node monitors the energy status of neighbours and triggers topology construction as needed. The framework of topology control approaches as shown in figure below:

![Fig. 1. Taxonomy of Topology Control][5]

In homogeneous topology control, the transmitting range of the nodes is the critical transmitting range that produces communicating graphs that are connected with high probability. Determining the critical transmission range using homogeneous topology control has been considered analytically as well as practically.

In Nonhomogeneous approach the transmitting range of the nodes varies and is classified into three categories, depending on the type of information that is used to compute the topology. In location based approaches, exact node positions are known. In direction-based approaches, it is assumed that nodes do not know their position, but they can estimate the relative direction of each of their neighbors. Finally, in neighbor-based techniques, nodes are assumed to know only the ID of the neighbors and are able to order them according to some criterion.

3. PROPOSED SCHEME

In this section we briefly describe the framework of ENR algorithm and its operating phases i.e. (1) topology construction phase and (2) topology maintenance phase.

3.1 Frame Work

The frame work of the ENR algorithm is described in the following steps as follows and as shown below in the figure 2:

Step-1: Initially, we get the energy of the randomly deployed nodes by using node energy model [1] that can accurately reveal the energy associated with each node.
Step-2: Next we will be determining the neighbors [2] of the sensor nodes; our method relies on distance estimation and the energy of the node.
Step-3: The maintenance of the proposed system is ensured by selecting the other node from the neighbor data structure [4] in case of energy depletion.
Step-4: Maintenance process involves periodic restoration and recreation of network topology. Here the each node periodically broadcasts the hello messages to its neighbor and waits for the acknowledgement.
3.2 Topology Construction

In topology construction phase, an efficient methodology is used that constructs a topology dynamically without violating network scenarios and thereby providing both connectivity and coverage. Hence a topology construction algorithm is proposed for building a topology and the steps for the same as shown in the figure 3.

We start constructing the topology \( T(v) \) where each node ‘n’ selects its neighbor ‘v’ based on range, energy and distance. Initially the neighbor list of a node is set to Null (line 2). In order to select the neighbor we determine the energy of each node \( E(v) \)(line 3-4). Now the node list \( E(v) \) is sorted in decreasing order(Line 6).After computing the energy we find the distance of each node ‘v’ from ‘n’ in order to determine the neighbor list of every node (Line 10-14).The nodes which are not present in the neighbor list \( N(v) \) is set to inactive state.

3.3 Topology Maintenance

In topology maintenance phase, as node start communicating with each other the energy depletes and some nodes fail rendering the network useless. Hence to make network functioning optimal we need to provide topology maintenance [4]. This process involves periodic restoration, rotation and recreation of network topology. Here the node periodically broadcasts the hello messages to its neighbor and waits for the acknowledgement(line 1-2).The Hello message carries the following information:<node_id,neighbor_id,energy>. If the node does not receive back acknowledgment (if the neighboring node is not having the energy to broadcast a message) then it invokes topology construction function. the steps as shown in figure 4.
4 Simulation Environment

4.1 Simulation model and parameters

We use ns2.33[7] to simulate our proposed protocol. In our simulation, we use the DCF of IEEE 802.11 for wireless LANs as the MAC layer protocol. It has the functionality to notify the network layer about link breakage.

Our simulation settings and parameters are summarized in table below:
<table>
<thead>
<tr>
<th>Parameter</th>
<th>Value</th>
</tr>
</thead>
<tbody>
<tr>
<td>Number of simulated Nodes</td>
<td>70</td>
</tr>
<tr>
<td>Pause time</td>
<td>10 sec</td>
</tr>
<tr>
<td>Area size of topography (m)</td>
<td>1600*2550 m</td>
</tr>
<tr>
<td>Radio range</td>
<td>250m</td>
</tr>
<tr>
<td>Packet size</td>
<td>512 bytes</td>
</tr>
<tr>
<td>Send rate of traffic</td>
<td>10 packets/sec</td>
</tr>
<tr>
<td>Traffic type</td>
<td>TCP</td>
</tr>
<tr>
<td>MAC Type</td>
<td>802.11</td>
</tr>
<tr>
<td>Simulation Time</td>
<td>150 s</td>
</tr>
<tr>
<td>Simulated Routing Protocols</td>
<td>ENR,CBTC,K-Neighbor</td>
</tr>
</tbody>
</table>

Table-1: Simulation Parameters.

4.2 Performance Metrics

We evaluate mainly the performance of ENR algorithm according to the following metrics, by varying the pause time.
Throughput: It is the number of bits that are conveyed or processed per unit of time.
Packet Delivery Ratio: It is the ratio of the no. of packets received successfully and the total no. of packets sent.
Bit Error Rate: It is the percentage of bits with errors divided by the total no. of bits transmitter, received or processed over a given time period.

4.3 Results

In this experiment, we measure the performance of the protocols at different pause time.
In Fig 5, Initially throughput increases with the increase in time. Initial increase in throughput is due to multiple route capability of ENR protocol (when the time increases from 0 to 5 sec). However as the time further increases the throughput decreases due to the dropping of packets due to the network congestion (this is shown when time increases from 5 to 50). Further from above plot we can determine our ENR protocol provides better throughput than CBTC and K-Neighbor protocol and maximum throughput is achieved when time is 5 sec.
In Fig 6, as the time increases, the number of possible paths between source and sink pair increases, which causes a significant increase in the exploratory messages within the network, initially leading to a higher packet delivery ratio of about 95-100% in the case of ENR. With further increases in time, the performance of packet delivery ratio decreases due to congestion, which in turn leads to collisions. However, in the case of ENR protocol, the packet delivery ratio remains somewhat stabilized as compared with CBTC and K-Neighbor.
In fig 7, initially when we are sending data the percentage of error rate is high due to noise or interference. However as the time increases the error rate of ENR is comparatively less as compared with CBTC and K-neighbor. As Error rate decreases in ENR the throughput increases, the error rate is decreased with the help of varying the FEC code size so that the erroneous packets are corrected instead of going for retransmission which improves the throughput.

5 Conclusion

In this paper we have proposed and implemented ENR (energy neighbour range) algorithm. This algorithm performs better in variable sized networks as compared with the the state of art methods. Our algorithm adopts a non-homogeneous approach which takes into consideration location, direction and neighbor information of a system in a network. However our algorithm results in some overhead as it needs to compute various parameters like energy, range and neighbors. This is overshadowed by high coverage and enhancing the lifetime of a network.

6 REFERENCES


AUTHORS

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