# DOUBLE FOUR-BAR CRANK-SLIDER MECHANISM DYNAMIC BALANCING BY META-HEURISTIC ALGORITHMS

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#### **ABSTRACT**

In this paper, a new method for dynamic balancing of double four-bar crank slider mechanism by metaheuristic-based optimization algorithms is proposed. For this purpose, a proper objective function which is necessary for balancing of this mechanism and corresponding constraints has been obtained by dynamic modeling of the mechanism. Then PSO, ABC, BGA and HGAPSO algorithms have been applied for minimizing the defined cost function in optimization step. The optimization results have been studied completely by extracting the cost function, fitness, convergence speed and runtime values of applied algorithms. It has been shown that PSO and ABC are more efficient than BGA and HGAPSO in terms of convergence speed and result quality. Also, a laboratory scale experimental doublefour-bar crank-slider mechanism was provided for validating the proposed balancing method practically.

### Keywords

Mechanism; balancing; BGA; PSO algorithm; ABC algorithm; HGAPSO algorithm

## **1. INTRODUCTION**

Dynamic balancing is one of the important concerns in high speed mechanisms such as crank slider mechanism, a special attention should paid to the inertia-induced force (shaking force) and moment (shaking moment) transmitted to the frame in order to reduce vibration amplitude. The methods of balancing linkages are well developed and documented in [1]. These techniques mostly are based on mass redistribution, addition of counterweights to the moving links, and attachment of rotating disks or duplication of the linkages [2]. In these methods, the shaking forces and shaking moments should be minimized. One of these methods is "Maximum recursive dynamic algorithm" that published by Chaudhary and Saha [3]. Another method which is documented by QI and Pennestrlis [4] is called "refined algorithm". It presents a numerically efficient technique for the optimum balancing of linkages. In this approach, instead of solving directly the dynamic equations, a technique is introduced to solve the linked dynamic equations in a "shoe string" fashion. The comprehensive mass distribution method, for an optimum balancing of the shaking force and shaking moment is used by Yu [5] to optimal balancing of the spatial RSSR mechanism.

In recent years, optimization techniques have received widespread interest in the engineering sciences. These techniques were initiated, primarily, by the advent of high speed digital computers which give a practical, economical and accurate means of obtaining solutions to

meaningful problems in the optimization of static and dynamic systems. Dynamic balancing inherently constitutes an optimization problem. For instance, CWB (counterweight balancing) involves a trade-off between minimizing the different dynamic re-actions. Therefore, determining the counterweights' mass parameters is an optimization problem. Also, Alici and Shirinzadeh [6] considered "Sensitivity analysis". In this technique they formulated the dynamic balancing as an optimization problem such that while the shaking force balancing is accomplished through analytically obtained balancing constraints, an objective function based on the sensitivity analysis of shaking moment with respect to the position, velocity and acceleration of the links is used to minimize the shaking moment. Evolutionary algorithms such as GA and PSO are useful in optimization problems, because they are easy to implement for the complex real-world problems. In some cases that objective function could not expressed as an explicit function of design variables, using these evolutionary algorithms is more suitable for optimization than the traditional deterministic optimization methods. Wen-Yi Lin proposed a GA-DE hybrid algorithm for application to path synthesis of a four-bar linkage [7]. Acharyya and Mandal[8]applied BGA with multipoint crossover, the PSO with the CFA (constriction factor approach) and the DE to the path synthesis of a four-bar linkage. Jaamiolahmadi and Farmani [9] proposed another related work. They described the application of GA to force and moment-balance of a four bar linkage. The aim of present work is applying new artificial intelligent based optimization methods for balancing double four-bar crank slider mechanism which is a special case of four-bar mechanism. but it has its own specification in dynamic balancing. Ettefagh and Abbasidoust [10] applied BGA for balancing of a crank slider mechanism, however in present work a comprehensive study of different algorithms proposed for balancing of a crank slider mechanism was carried out. The major objective of this paper is to compare the computational effectiveness and efficiency of BGA (Binary Genetic Algorithm), PSO (Particle Swarm Optimization), ABC (Artificial Bee Colony) and HGAPSO (Hybrid GA and PSO) algorithms. The performance of the optimization techniques in terms of computational time and convergence rate is compared in the results section. After kinematics formulating of the mechanism and extracting the object function with its' constraints, results of the proposed method on a simulated mechanism will be reported. Finally, for experimentally validation of the method, a laboratory scale experimental crank slider mechanism will be considered and the proposed optimization method will be applied for its balancing.

## **2. MODELLING**

The selected mechanism in this paper is shown in Figure 1, which is a benchmark double four-bar mechanism and this mechanism is applied in different machines [10]. Therefore the proposed balancing method in this paper may be easily applied to other mechanisms in different machines for balancing. As shown in Figure 1, the mechanism contains two cranks-sliders and four rotating disks. Each rode is attached to disk 1 and disk 4. All disks are connected to a shaft which is rotating with an electromotor. When the mechanism is working, shaking force and moments cause frame to vibrate. Our goal is to minimize shaking forces and moments according to the regarded constraints. The mechanism is modeled as shown in Figure 2 by computing shaking forces of the mechanism on all four planes.



Figure 1.Schematic diagram of the considered double four-bar crank-slider mechanism [10]



Figure 2.Shaking forces of the mechanism on all four plans [10]

In this paper, the method used for balancing is the addition of two counterweights on disks 2 and 3. In addition, system has an unbalance mass on disk 2. Shaking forces in crank slider as the dynamical forces of mechanism are formulated as:

$$f_c = m_c R \omega^2 \tag{1}$$

$$f_p = m_p R \omega^2 \left(\cos\theta + \frac{R}{L}\cos 2\theta\right)$$
<sup>(2)</sup>

Where  $m_c$  is the eccentric mass and  $m_p$  is the equal mass of the slider, *R* is corresponding radius and  $\omega$  is mechanism rotating speed. Also,  $\theta$  indicates the angular position of the disk and *L* is the connecting rod's length. Equilibrium on *x* and *y* axes due to the acting forces are described by:

$$\sum F_{x} = 0 \rightarrow$$

$$m_{p}R\omega^{2}(\cos\theta + \frac{R}{L}\cos 2\theta) + m_{c}R\omega^{2}\cos\theta + m_{0}R_{0}\omega^{2}\cos(\theta + \alpha) + m_{1}r_{1}\omega^{2}\cos(\theta + \varphi_{1}) + m_{2}r_{2}\omega^{2}\cos(\theta + \varphi_{2}) +$$

$$m_{p}R\omega^{2}(\cos(\theta + \theta_{0}) + \frac{R}{L}\cos 2(\theta + \theta_{0})) + m_{c}R\omega^{2}\cos(\theta + \theta_{0}) = 0$$

$$\sum F_{y} = 0 \rightarrow$$

$$m_{c}R\omega^{2}\sin\theta + m_{0}R_{0}\omega^{2}\sin(\theta + \alpha) + m_{1}r_{1}\omega^{2}\sin(\theta + \varphi_{1}) + m_{2}r_{2}\omega^{2}\sin(\theta + \varphi_{2}) +$$

$$m_{c}R\omega^{2}\sin(\theta + \theta_{0}) = 0$$
(3)

Also Momentum equilibrium equations are written as:

$$\sum M_{y} = 0 \rightarrow [m_{0}R_{0}\omega^{2}\cos(\theta + \alpha) + m_{1}r_{1}\omega^{2}\cos(\theta + \varphi_{1})]a_{1} + [m_{2}r_{2}\omega^{2}\cos(\theta + \varphi_{2})]a_{1} + a_{2}) + [m_{p}R\omega^{2}(\cos(\theta + \theta_{0}) + \frac{R}{L}\cos(2(\theta + \theta_{0}))](2a_{1} + a_{2}) + [m_{c}R\omega^{2}\cos(\theta + \theta_{0})](2a_{1} + a_{2}) = 0$$

$$\sum M_{x} = 0 \rightarrow [m_{0}R_{0}\omega^{2}\sin(\theta + \alpha) + m_{1}r_{1}\omega^{2}\sin(\theta + \varphi_{1})]a_{1} + [m_{2}r_{2}\omega^{2}\sin(\theta + \varphi_{2})](a_{1} + a_{2}) + [m_{c}R\omega^{2}\sin(\theta + \theta_{0})](2a_{1} + a_{2}) = 0$$
(5)

 $m_0$  is the constrain mass on disk 2,  $R_0$  and  $\alpha$  are corresponding radius and angular position of the mass, respectively.

In construction of optimization problems, as described before, some parameters of system is applies as the variables of cost function to minimize it. In our method,  $m_1$ ,  $m_2$ ,  $\varphi_1$ ,  $\varphi_2$ ,  $r_1$  and  $r_2$  are counterweights parameters that are unknown. We take  $r_1$  and  $r_2$  as known parameters. So we define the cost function as below. The cost function is consisting of forces equations and the constraints are considered by moment equations. For this purpose, we define some parameters as  $P_1$ ,  $P_2$ ,  $P_3$  and  $P_4$  as below:

$$P_1(\theta) = \sum F_{X} \quad , \ P_2(\theta) = \sum F_{Y} \tag{7}$$

$$P_3(\theta) = \sum M_x \, , \, P_4(\theta) = \sum M_y \tag{8}$$

These parameters are the sum of forces and moments in x, y direction. Consequently, the cost function is:

Cost function: 
$$f(m_1, m_2, \varphi_1, \varphi_2) = \int \{ |P_1(\theta)| + |P_2(\theta)| \} dA$$
 (9)

The cost function is a function of the parameters m1, m2,  $\varphi 1$  and  $\varphi 2$ , which should be found. In other words, the cost function is the area that  $\{|P_1(\theta)| + |P_2(\theta)|\}$  makes in polar system. Our goal is to minimize this area. Additionally, by defining C1 and C2 as follows:

$$C1(m_1, m_2, \varphi_1, \varphi_2) = \int |P_3(\theta)| dA$$
(10)

$$C2(m_1, m_2, \varphi_1, \varphi_2) = \int |P_4(\theta)| dA \tag{11}$$

Constraints are considered such that C1 and C2 to be less than predefined values which determined by searching among the maximum value of them during running the Meta-Heuristic Algorithms.

## **3.** APPLYING META-HEURISTIC ALGORITHMS

PSO, ABC and HGAPSO algorithms, which is introduced in summary in following section, are applied for minimizing the mentioned cost function with defined constraints as it was described in previous section.

### 3.1. Particle Swarm Optimization algorithm

The PSO method introduced for the first time in 1995 by Kennedy and Eberhart [11]. This method based on social behavior of organisms such as fish schooling and bird flocking. Because of its simple concept and quick convergence, PSO can be applied to various applications in different fields. The approach uses the concept of population and a measure of performance like the fitness value used with other evolutionary algorithms. Similar to other evolutionary algorithms, the system is initialized with a population of random solutions, called particles.

Each particle maintains its own current position, its present velocity and its personal best position explored so far. The swarm is also aware of the global best position achieved by all its members. The iterative appliance of update rules leads to stochastic manipulation of velocities and flying courses. During the process of optimization the particles explore the *D*-dimensional space, whereas their trajectories can probably depend both on their personal experiences, on those of their neighbors and the whole swarm, respectively. In every iteration, each particle is updated by following the best solution of current particle achieved so far (*Pbest*) and the best of the population (*gbest*). When a particle takes part of the population as its topological neighbours, the best value is a local best. The particles tend to move to good areas in the search space by the information spreading to the swarm. The particle is moved to a new position calculated by the velocity updated at each time step t. This new position is calculated as the sum of the previous position and the new velocity by equation (12).

$$v_i^{t+1} = wv_i^t + c_i rand(0,1) (Pbest_i - x_i^t) + c_i rand(0,1) (gbest_t - x_i^t)$$
(12)  
The velocity updated as following equation:  
$$x_i^{t+1} = x_i^t + v_i^{t+1}$$
(13)

The parameter w is called the inertia weight and controls the magnitude of the old velocity 
$$v_i^t$$
 in

the calculation of the new velocity which  $isv_i^{t+1}$ , whereas the acceleration factors C<sub>1</sub> and C<sub>2</sub>determine the significance of *Pbest* and *gbest* respectively. In the present work we selected C<sub>1</sub>=0.25 and C<sub>2</sub> =0.15. Furthermore,  $x_i^t$  is the current particle position and  $x_i^{t+1}$  is the new position of the particle. *rand*(0,1) is the random number uniformly distributed. In general, the inertia weight *w* is set according to the following equation:

$$w = w_{max} - \left(\frac{w_{max} - w_{min}}{lr_{max}}\right) \times Iter$$
<sup>(14)</sup>

Where  $w_{max}$  is initial weight,  $w_{min}$  is final weight,  $Ir_{max}$  is the maximum number of iteration (generation) and *Iter* is current iteration number.

Main steps of the procedure are:

- 1. Position and initial velocity of particles are generated randomly.
- 2. The objective function value i calculated for each particle.

- 3. *Pbest* For each particle is considered equal to its initial position (In the first iteration). Also *gbest* determined.
- 4. Modify position and velocity of each particle using the equations (12) and (13).
- 5. Calculate the overall objective function for each particle.
- 6. If the value of new overall objective function is better than the value of each particle in *P*-*best*, then it would be replaced in *P*-*best* and if in whole new population gets found which is better fitting in objective function would be replaced by *g*-*best*.
- 7. Repeat steps 4 to 7 to reach the maximum number specified .The final answer is obtained from the *gbest* in the last iteration.

Figure 3 shows the flowchart of the developed PSO algorithm [11].



Figure 3. Flowchart of the PSO algorithm

#### 3.2. Artificial Bee Colony algorithm

The artificial bee colony (ABC) algorithm is a new population-based stochastic heuristic approach which was introduced for the first time by Karaboga in 2005 [12]. This algorithm is very simple and flexible, which does not require external parameters like crossover rates, especially suitable for engineering application. The colony of artificial bees consists of three groups of bees to search foods generally, which includes employed bees, onlookers and scouts. The first half of colony consists of the employed artificial bees and the second half includes the onlookers. The employed bees search the food around the food source in their memory. They perform waggle dance upon returning to the hive to pass their food information to the other of the colony (the onlookers). The onlookers are waiting around the dance floor to choose any of the employed bees to follow based on the nectar amount information shared by the employed bees. The employed bees is equal to the number of food sources around the hive. The onlooker bee chooses probabilistically a food source depending on the amount of nectar shown by each employed bee, see equation (15).

$$P_i = \frac{fit_i}{\sum_{n=1}^{SN} fit_i} \tag{15}$$

Where  $fit_i$  is the fitness value of the solution *i* and *SN* is the number of food sources which are equal to the number of employed bees.

Each bee searches for new neighbor food source near of their hive. After that, employed bee compares the food source against the old one using equation 16. Then, it saves in their memory the best food source [13-14].

$$V_{ij}^{new} = x_{ij}^{old} + \varphi_{ij} \left( x_{ij}^{old} - x_{kj}^{old} \right)$$
(16)

Where  $k \in \{1,2,3,...,NP\}$  and  $j \in \{1,2,3,...,n\}$  are chosen randomly.  $\varphi_{ij}$  is a random number between [-1, 1]. In ABC, Providing that a position cannot be improved further through a predetermined number of cycles, then that food source is assumed to be abandoned. The value of predetermined number of cycles is an important control parameter of the ABC algorithm, which is called "limit" for abandonment. Assume that the abandoned source is  $x_i$  and  $j = \{1, 2..., N\}$  then the scout discovers a new food source to be replaced with $x_i$ . This operation can be defined using equation 17.

$$x_{i}^{j} = x_{j}^{\min} + \operatorname{rand}(0,1)(x_{j}^{\max} - x_{j}^{\min})$$
 (17)

ABC implementation stages of the algorithm:

- 1. Initialization
- 2. Place employed bees on the food sources
- 3. Place onlookers on the food sources
- 4. Send scout to carry out the global search

Flowchart of the ABC algorithm is shown in figure4 [12].



Figure 4. Flowchart of the ABC algorithm

## **3.3.** Hybrid of Genetic Algorithm and Particle Swarm Optimization (HGAPSO)

#### Concept of GA:

Genetic algorithm was introduced by John Holland in 1967. Later, in1989, by Goldberg's attempts got accepted. And today, due to its capability, it has prosperous place among other methods. Routine optimization in genetic algorithm bases on random- directed procedure. In this way depending on Darwin's theory of evolution and fundamental ideas the method has been firmed. First of all, for some constants which are called population, a swarm of main parameters, eventually get produced. After performing scalar simulator program, a number that is attributed as standard deviation or set fitness data, impute as a member of that population. Fitness process is done for all created members then by summoning genetic algorithm operators such as succession (crossover), mutation and next generation selection (selection) are formed. This routine continues till satisfying convergence [15].

#### Concepts of PSO:

Initially, velocity and position of each particle randomly select and direct all particles gradually to the best answer. PSO has high speed in local and total search. Applying PSO takes to dominate rapid and untimely convergence of GA that is achieved with elitism operations. Also, PSO can fit particles of each generation whereas GA is seeking for an answer generation to generation. Depending on these points, GA and PSO are good supplements of each other [16].

The HGAPSO algorithm is following as bellow:

- 1. A population randomly should be selected, initially, which is parameters of GA and particles of PSO.
- 2. The fitness of concluding is calculating.
- 3. Half of population with high fitness would be selected and then PSO function would be performed. But a factor breeding ratio was presented with  $\varphi$ symbol [17]. The factor determines the percentage of population on which the PSO operation is done.
- 4. Enhanced elites would be transmitted to the next generation directly and mutation and remixing functions would be done on the rest population.

Figure 5 shows all the above steps and flowchart of the HGAPSO algorithm is shown in Figure6 [16].



Figure 5. Steps of HGAPSO algorithm [16]



Figure 6. Flowchart of the HGAPSO algorithm

## 4. RESULTS AND DISCUSSION

In this section the results of mentioned algorithms implementation in our optimization problem for minimizing cost function will be presented and discussed. Table 1 expresses the results of dynamic balancing optimization with PSO, ABC, BGA and HGAPSO algorithms for 200 and 300 iterations. Also programming is done with MATLAB12, with 2.8 GHZ computer system processor and 1024MB memory. By these tables, best and worst average results are shown in bold. Results comparison indicates that PSO algorithm is more efficient than other algorithms, as the average cost of objective function for 20 iterations is 29853.35062, and for 300 iterations is 29853.30955 that is less than three other algorithms.

Table 1. Results of dynamic balancing optimization with PSO, ABC, BGA and HGAPSO algorithms for
200 and 300 iterations

Statistical Results Obtained by the BGA algorithm for 200 Generation         Experiment         M1         M2         Fit         Fi2           #1         425.2199         770.088         2.200227         4.907356         2.99177403         27.81430         #1         967.917         87.0169         9.831981         4.470459           #2         200         312.61         3.14464         4.701646         3.00277         2.9852082         2.8.21068         #2         4.68.1223         857.792         8.58609         5.188800           #3         425.2199         648.6004         2.395364         4.54247         2.9982.0925         2.8.21068         #4         2.21.443         1.07.222         1.19885           #4         305.5718         3.74.105         3.722004         5.200833         2.9844086         2.980209         8.80209         6.23.1133         8.40666         11.19885           #6         38.8517         7.99.444         1.228364         4.391474         2.9928.107         2.070006         #7         12.15862         10.31822         6.404841           #7         1.181.5777         1.181.5777         1.9912797         1.2014977         1.983979         8.803097         8.803097         8.803097         8.803097         8.803	Cost 29853.8571 29854.7679 29856.3607 29854.8753 29855.1337 29854.5122	CPU Time 24.515309
Statistical results Oblatified by the bick algorithm for 200 Generation         CPU Time         Mi         M2         Fit         Fi2           Experiment         Mi         M2         Fit         Fit         CPU Time         Fit	Cost 29853.8571 29854.7879 29856.3607 29854.8753 29855.1337 29855.1337	24.515309
Experiment         M1         M2         Fi1         Fi2         Cost         CPU Time         Fi1         20/2017         40/2019         420/000         420/000           #1         425.2199         770.08         2.20027         49/0736         29/31/201         21.81440         47.0180         20.81640         448.3323         458.3323         857.992         8.58098         5.88482           #2         200         312.81         3.14464         4.710854         30037.9671         29.93368         #4         1131.1942         1101.4856         9.255999         1.8518603           #4         305.5716         3.71038         2.92046         2.20048         2.852836         4.84066         11.1989           #5         312.81         648.604         2.352366         4.870685         2.9914.4407         2.985386         #6         677.15         2.21.4425         10.7222         1.12689           #6         398.515         798.4426         1.2923.84         2.9923.807         3.2070005         #6         677.15         2.21.5824         10.317827         6.045841           #7         1131.177         1131.177         1007.107         2.9923.8173         2.9070056         #7         12.78584         90.86411 <td>29853.8371 29854.7879 29856.3607 29854.8753 29855.1337 29854.5122</td> <td>24.010309</td>	29853.8371 29854.7879 29856.3607 29854.8753 29855.1337 29854.5122	24.010309
42         100         2.00221         4.00224         4.00254         2.000         2.00025         2.000         2.00000         2.00000         2.00000         2.00000         2.000000         2.000000         2.000000         2.000000         2.000000         2.000000         2.000000         2.000000         2.00000000         2.0000000         2.0	29856.3607 29854.8753 29855.1337 29854.5122	24 497529
#3         425.2199         648.6004         2.3953.80         4.94.4247         29982.0205         28.827008         #4         217.0072         221.443         107.222         11.26335           #4         306.5718         37.1055         3.722004         6.2200633         25844.0966         28.86260         #5         260.4993         6.23133         8.480066         11.19885           #6         305.515         778.4428         1.223884         4.391474         22.9828.107         22.078005         #6         677.15         327.1582         10.3722         6.49843           #6         305.515         778.4428         1.223884         4.391474         29.928.107         22.078005         #7         12.26584         900.89413         10.015823         8.862907           #7         1131.0777         10.91097         27.91787         20.911497         72.912868         #8         90.89413         10.015823         8.862907	29854.8753 29855.1337 29854.5122	25.638495
#4         305 5718         374 1535         3.722004         5.220833         298540386         28.386290         #5         200.4933         6.23.3153         8.486066         11.19855           #5         312.61         646.8604         2.352365         4.876685         29914.4407         32.985386         #6         677.15         321.582         10.37823         6.049343           #6         368.915         789.4282         12.23384         4.391474         29928.107         32.076005         #7         1237.854         960.86413         10.19323         6.862037           #7         1131.5777         10.914007         73.273753         200.164.14         78.1267         748.3797         748.3797         5454591	29855.1337 29854.5122	25.814662
#5         312.61         648.6804         2.352358         4.876685         23914.4407         32.985386         #6         677.15         321.5862         10.37822         6.049843           #6         589.515         789.4428         1.232354         4.391474         29928.107         32.078005         #7         122.8654         960.86413         10.019223         5.862087           #7         1131.5777         1010.14100         75.77375         300.15414         29.154391         5.945891	29854.5122	26.448376
#6 368.915 789.4428 1.228384 4.391474 2.9928.107 32.076005 #7 1237.6854 960.84413 10.019323 5.8620878 #7 1237.6854 960.84413 10.019323 5.8620878		28.761479
#7 1183 5777 1099 1202 2 9419802 5 7979735 30043 6419 28 184769 #8 746.1877 749.8799 6.339937 9.548391	29856.5621	27.566192
	29854.274	27.581421
#8 555.4252 310.8504 3.924688 6.04365 29876.8727 28.312547 #9 6.39.291/4 315.52(13 7.741611 12.860/6	29855.4856	27.869337
#9 425 2199 873,5003 13932/73 4,802392 23961/9256 315,78058 **** 600,500 000,000 27,7242* 0540000	29855.04881	26.6315405
TIU 536.0/04 40/.6246 5.312/62 3.136522 25957.1166 30.2(1240 Best Best Best Best Best Best Best Best	29853.8571	24.497529
Best 22854.1 27.81643 Worst	29856.5621	28.761479
Statistical Results Obtained by the BGA algorithm for 300 Generation           Experiment         M1         M2         Fit         Fiz         Cent         CPU Time	ithm for 300 I	Iterations
#1 585.3372 416.4223 3.777281 5.94538 29860.512 69.617549 #1 889.97288 1026.2293 6.8552951 13.822746	29856.0266	40.422035
#2         537 8299         845.7478         0.393083         3.881694         29979.2346         75.797888         #2         374.24233         1405.1666         7.839461         10.87474	29859.7044	42.796960
#3 363.6384 536.0704 3.040251 5.208349 29900.57 70.299023 #3 260.895 501.8796 7.752119 10.27209	29854.4018	40.764053
#4 6205279 0.7861659 4.121229 256.305 29988.9287 61.185702 #4 1624.3941 1153.022 10.633852 7.5664349	29855.7171	43.498866
#5 397.0674 648.6804 2.745439 5.128504 29896.1839 67.436636 #5 843.7006 359.6281 11.00391 8.106481	29854.1593	45.076684
#6 321.4076 529.0323 5.920812 3.90012 2.9866.535 7.3210396 #6 772.2893 653.4779 5.531582 9.657654	29854.0499	48.079070
#7 223.1525 590.9091 1.94064/ 4.710654 2966.4003 85.515155 #7 253.2995 1494.202 4.5137612 7.5742005 #8 200 564.2923 1.67060 4.46577 2926.7044 6.7214300 #8 1.401.075 1182.7054 10.982768 0.91466688	29854.0059	48.203408
#9 256.6217 310.5514 3.24588 5.67055 298417926 65.565551 #9 722.1371 560.479 9.61366 6.1950000	25054.0752	49.274214
	29855 4519	
#10 485.044 423.4604 3.531605 5.560567 2.9900.5253 77.98145 #10 655.3966 612.1464 9.454685 5.518574	29855.4519 29853.7587	45.614997
#10         485.04         423.4604         3.531605         5.65067         2990.5253         77.9511.15         #10         655.3656         612.1464         9.454685         5.518574           Average          2.2999.98         71.41247         Average         Average          5.518574	29855.4519 29853.7587 <b>29855.22148</b>	45.614997 45.1617406
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Also, above mentioned results in the table are illustrated in Figures7 and 8 in bar chart which show the comparison of objective function cost. By observing these figures, it may be concluded that the cost order is as following: PSO<ABC<BGA<HGAPSO.



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Figure 7. Comparing cost function value for 200 Iterations



Figure8. Comparing cost function value for 300 Iterations

In addition, the runtime comparison shows that PSO algorithm finds the optimal solution in shorter time than the others. Also the average runtime for this algorithm for 200 iterations is 13.2666839 seconds and for 300 iterations is 21.4484653 seconds. According to numerical results PSO algorithm is 3.84 times faster than BGA and 2.56 times faster than HGAPSO and 2.21 times faster than ABC. Figures 9 and 10 indicate the runtime of these algorithms for one performance with 200 and 300 iterations, respectively. By observing these figures, it maybe found that the runtime order is as following: PSO<ABC<HGAPSO<BGA. It should be mentioned that the runtime for Hybrid algorithm of BGA and PSO algorithm is 1.5 times faster than BGA algorithm in reaching optimal solution.



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Figure 9. Comparing runtime for 200 iterations



Figure 10. Comparing runtime for 300 iterations

Figure 11 indicates the time order of these algorithms for one performance with 100 and 200 repetitions. In this chart, the implementation of time progress of these algorithms is in term of seconds.



Figure 11. Computation time variation versus generation for 100 iterations (Left) and 200 iterations (right)

Figure 12 shows fitness functions of four algorithms in 100 runs. In this figure the range of fluctuation of fitness value for BGA and HGAPSO algorithms is between 29853 and 29855, while this range for BGA and HGAPSO is approximately between 29853 and 30200. Also high difference between the best and worst result of the BGA and HGAPSO algorithms indicates the weakness and dependence of these two algorithms and their sensitivity to the initial population, while PSO and ABC solution range is around the mean values, so it indicates the stability of them. Also, the PSO and ABC algorithms are close to average result that this matter indicates the stability of these two algorithms.



Figure 12. Fitness functions fluctuation for 100 runs for four algorithms

Figure 13 and 14 display the four algorithms convergence in a sample run of them (for 200 and 300 iterations). According to the minimization target, an important point which should be considered is that the charts are descending. In addition the scale of  $10^4$  is using for measuring of objective function. As comparing of these two figures, it may be observed that the PSO and ABC algorithms get convergence approximately after 30 generations also they achieve the optimized solution in high speed rather than HGAPSO and GA.



Figure13.Convergence plot for 200 iterations



Figure 14. Convergence plot for 300 iterations

Finally, Figure 15 indicates cost function diagram in balanced and unbalanced states by mentioned algorithms in polar plot. This type of plot is a standard plot which applied for illustrating the ability of the proposed method in balancing of the mechanism in one revolution of the mechanism driver (0-360°). As it was described in previous sections, the main purpose of optimization is minimizing of the cost function in equation 9 which maybe interpreted as minimizing the limited area in the mentioned plot. Therefore better balancing method is a method which has minimum and uniform area in the polar plot. By observing this figure, it maybe concluded that PSO has better performance than other algorithm, because it contains minimum area in comparing with other algorithms.



Figure 15. Polar plot of cost function for unbalanced and balanced states by applying different algorithms

## 5. EXPERIMENTALLY VALIDATION OF THE PROPOSED METHOD

For showing the practicability of the proposed balancing method based on mentioned optimization algorithms, a laboratory scale experimental double four-bar crank-slider mechanism was provided as shown in Figure 16 [10]. By using two accelerometers (type 4507 B&K Co.) and data acquisition system (PULSE B&K Co.), the acceleration of the right and left bearing of the mechanism can be extracted. Also a tachometer (Type 0024 MM B&K Co.) is applied for identifying the one revolution of the crank for estimating the angular acceleration. As a sample, the acceleration signals of the unbalanced and balanced mechanism, applying BGA, are illustrated in Figures 17 and 18. By comparing the mentioned linear and angular acceleration signal of the unbalanced and balanced mechanism, the possibility and accuracy of the proposed method in balancing of real mechanisms can be observed. In addition the corresponding angular acceleration signals are illustrated in Figure 19. As it is obvious from these figures, the proposed method is able to balance the mechanism because the overall amplitude of the linear and angular acceleration was decreased after balancing. Other similar results may be obtained by applying other algorithms, which cannot be illustrated in this paper because of the page limitation. By considering the experimental result, it can be observed that the same conclusion maybe derived as in previous section while comparing these algorithms.



Figure 16.A laboratory scale experimental double four-bar crank-slider mechanism setup



Figure 17. Acceleration signal of the right part of the mechanism for before and after balancing



Figure 18. Acceleration signal of the left part of the mechanism for before and after balancing



Figure 19. Angular acceleration signal of the mechanism for before and after balancing

## **6.** CONCLUSION

In this paper, four Meta-Heuristic algorithms, that are PSO, ABC, BGA and HGAPSO, were applied for balancing of double four-bar crank slider mechanism. This mechanism is a basic and benchmark mechanism which is applied in different machine and engines. Therefore proposed balancing method may be applied for different type of machines for reducing the vibration and noise because of unbalancing. Kinematic modelling of the mechanism and suitable objective function which should be minimized and the constraints are derived. In order to comparing the optimization algorithm, different results based on cost function, fitness, convergence speed and runtime values were extracted and studied. The results show that the PSO algorithm converges quickly to the optimum solution, while the optimality of PSO solutions has a higher quality than other algorithms. Also, reviewing the responses of the four algorithms in several performances with various primary populations determined that, PSO and ABC are more stable than BGA and HGAPSO. Therefore PSO and ABC have a good potential in finding optimal solution and its convergence characteristics are more favourable. In addition, the practicability of the method was proved by applying the proposed method on a laboratory scale experimental complex crank slider mechanism.

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