

EVALUATIONS OF USER CREATION PERSONAL PORTAL PAGE USING DACS WEB SERVICE

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ABSTRACT

A personal portal, which is an entrance wherein each user can acquire the information that s/he is interested in on a network, is often used as an alternative means of communication. However, there are a number of problems with the existing personal portals. For example, because the Web page as a personal portal is generated by the program located on the specific Web server which is managed by a system administrator, it is not always ideal for all users. To solve this kind of problems, we developed two Web Service functions, which are realized on the network by introducing the Destination Addressing Control System (DACS) Scheme. These two Web Service functions are as next. The first is the function to extract the data for each user from a database and display it on the Web browser. The second is the function to retrieve the data for each user from a document medium and display it on the Web browser. Through these Web Service functions, each user can easily create a customized personal portal that displays personal information. In this paper, the above two functions are extended to manage information not only for each user but also for each group of users and for all users, and the extended two functions are integrated as a DACS Web Service. By using the DACS Web Service, each user can create and customize the Web page as a personal portal for practical usage in an individual organization. After the prototype system's implementation, evaluations are performed.

KEYWORDS

Personal portal, Web Service function, DACS Web Service

1. INTRODUCTION

Static Web pages are often used as an alternative means of communication in addition to e-mail by an unspecified number of users. However, they are unsuitable for communication among individual users for their individual purposes or interests. As an alternative means of communication, the personal portal is realized by a Web Service; it allows each user to change the contents of the Web page dynamically. The term "Web Service" refers to a service that is provided to users through a Web server, and the term "portal" often indicates a Web page for searching information [1][2] such as with Google or Yahoo. The meaning of the phrase "personal portal" used in this paper is different from that of the above "portal" as just described. Here, "personal portal" refers to an entrance where each user can acquire information that s/he is interested in on the network; it can display different information for individual users dynamically on a Web browser by using a program such as the Common Gateway Interface (CGI) [3]. As examples, Netvibes [4] and Google [5] are introduced. Therefore, personal portals are suitable for communication among individual users for their respective purposes or interests. To display the necessary information for an individual user on a Web browser, the information is searched

and extracted from databases on the network, after which each user is notified by a program such as CGI. If each database is distributed on the network, the program used to retrieve the information is large and complicated. Given that the program is installed by a system administrator and cannot be freely changed by a user, using a Web page as a personal portal is not always ideal for all users. To solve this problem, a new form of personal portal, which each user can create and customize easily, is necessary. The Destination Addressing Control System (DACS) Web Service is proposed and examined to realize such a personal portal. This DACS Web Service is implemented by extending two types of Web Service functions, which are realized on the network introducing the DACS Scheme. Using this DACS Web Service, each user can easily create and customize a Web page as a personal portal for oneself.

The DACS Scheme is a network management scheme. The basic principle of the DACS Scheme is that the entire network system is managed through communication control for each user on the client machine [6]. Moreover, functional extension is achieved by allowing the communication control for each user to coexist with communication control for each client machine to use in a user group [7],[8]. In addition, Secure DACS Scheme is proposed to solve a security problem [9]. Two types of Web Service functions are primitive functions as fundamental elements of DACS Web Service, which are also realized on the network introducing the DACS Scheme. However, they have following two functions. In the first function, data that is stored in the database and is dispersed on the network can be used efficiently [10]. In the second function, data that is stored in a document medium such as PDF or simple text files can be used efficiently. Using both functions, when different users input the same URL into the Web browser, the different information for each user is searched and extracted from the database or document medium and displayed on the Web browser. By incorporating various kinds of URLs into a static HTML, it has become possible for each user to create and customize a Web page as a personal portal easily [11]. However, in some user groups, the function of sending information not only to a specific user but also to a group of users or to all users is also necessary. In this paper, by extending and integrating two types of Web Service functions, the DACS Web Service is realized [12]. The DACS Web Service uses data stored in a database or document medium to enable individual users, groups of users, and for all users to send and receive information through the integrated user interface. This paper extends prior work, portions of which were presented in a conference paper [12]. The following points have been added from the conference paper.

(1) Revision of the existing studies and technology

In the conference version, only the basic mechanism of the personal portal was described. However, to clarify the differences with existing researches and technologies, more detail comparisons with them are performed by using additional figures (Fig.2 and Fig.3).

(2) Addition of the VPN communications

Because communications between the client and the DACS Web Service were not VPN communication in the conference version, the VPN communications are introduced to keep security in the journal version, that is explained by using the additional figure (Fig.5).

(3) Revision of the content of the DACS Web Service

In the conference version, summary of the DACS Web Service and two functions to constitute it (Function α and Function β) were described. To have readers understand the DACS Web Service correctly, revisions are performed overall. First, as Section IV A, operational procedures for personal portal are added newly to have readers understand the flow for creating the Web page of the personal portal. In this part, the method that users know where data exists is also presented by

using the additional figure (Fig.12). In Section IV B, two functions (Function α and Function β) that constitute the DACS Web Service are explained in detail by adding Fig.14. In addition, two extended functions of acquiring data for each group are explained more concretely and in detail in Fig.15 and Fig.16. As a result, the contents of Section IV C and Section IV D in the conference version are merged in Section 4.2 of the journal version. Then, Section IV C is added newly in the journal version to show the DACS Web Service's characteristics.

(4) Additional implementation of prototype system and evaluation

In the conference version, a simple Web page as a personal portal was described as experiment results. However, to have readers of this paper understand the contents of this paper more concretely, experiments by additional implementations of prototype system and evaluations are performed, and results of them are described in Section V of the journal version. As a result, Fig.18 is added newly, and the contents of Fig.19 and Fig.20 are changed.

2. RELATED WORKS AND MOTIVATION OF THIS STUDY

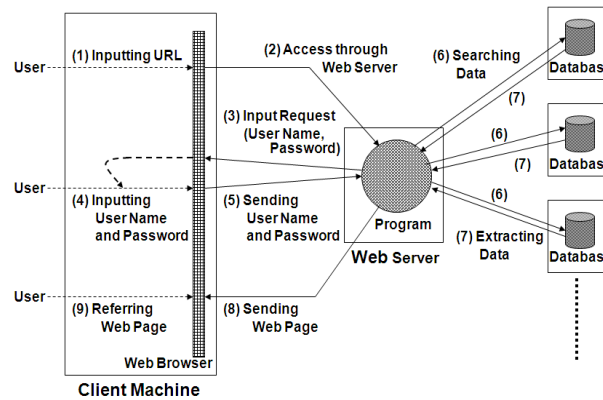


Fig. 1 Basic Mechanism of Existing Personal Portals

Existing personal portals are Web pages that serve as unified windows for data stored in the information systems of organizations. In the field of the personal portals, various studies have examined [13], [14], [15]. In addition, personal portals are developed in individual organizations for practical purposes, and commercial software packages are also used. The system known as “HIInT” is an example developed in a university [16]. “Blackboard” [17] is an example of a software package often used in a university. In these cases, the basic mechanism of personal portals is described by the sequence of processes from (1) to (9) shown in Fig. 1. The arrows with the dotted lines show the actions that the user performs in the Web browser, whereas those with solid lines show the flow of processing. This mechanism is explained in sequence as follows. First, the user enters a URL as input into the Web browser (1). Next, the Web server corresponding to the URL is accessed (2), and the program corresponding to the URL is executed. At this point in time, the program on the Web server side does not have the user information (user name and password). Therefore, an input demand for user information is executed on the Web browser side (3). Subsequently, the user enters the user information as input (4), which is sent to the program on the Web server (5). In existing mechanisms, user authentication is performed at this point in time. If access is permitted, information related to the user is searched from data that is accumulated in advance (6). To accumulate the data, there are methods using relational databases or document medium. Following search, the data related to each user are extracted (7). Programs on the Web server side programs that receive the data, such as CGI, generate a Web page dynamically from the data. Then, the Web page is sent to the Web browser side (8). Thus,

each user can view information which is related to him/herself and displayed in the form of a Web page (9).

Moreover, as one of the methods to realize the personal portal, there is a method by application frontend integration in Service Oriented Architecture (SOA) [18][19]. As shown in Fig. 2, SOA, which is the study of information system integration in organizations such as enterprises and universities, is realized by using WebAPIs. The mechanism is different from that of Fig. 1 in at least two points.

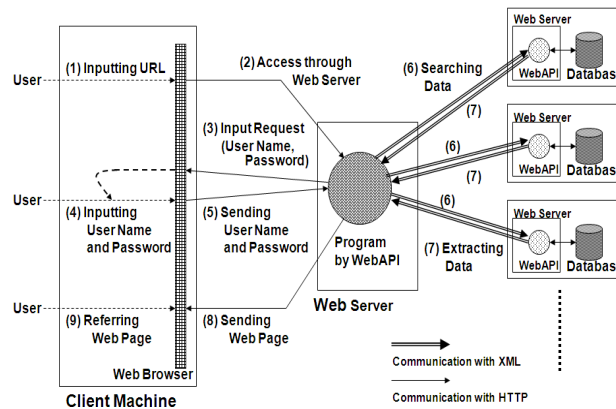


Fig. 2 Mechanism of Personal Portals by WebAPI

(1) The program on the Web server that the user accesses first is implemented with Web APIs such as Simple Access Object Protocol (SOAP) [20][21][22] and Representational State Transfer (REST) [23][24][25].

(2) Communication between (6) and (7) is performed with use of Extensible Markup Language (XML) [26][27] through each Web server distributed on the local area network.

The essential commonality of the mechanisms explained by here in Section II is that the program on the Web server extracts data from some database and generates the Web page as a personal portal. However, there are two problems with this general scheme.

(Problem 1) The program on the Web server is installed by a system administrator, and usually cannot be changed. Given that this program generates the Web page, it is not always easy for the user to customize the personal portal. Although this can be done in some cases within the specifications of the program, not all users can customize the Web page because the customization beyond these specifications is impossible.

(Problem 2) Although data from a database are used in existing personal portals, the data are often stored in the form of particular document medium in individual organizations. However, existing personal portals cannot use such data easily.

To solve these problems, two types of Web Service functions based on the DACS Scheme were proposed. By using the first function [10], problem 1 was solved. This enabled each user to create a Web page as a personal portal using data from the database. By using the second function [11], problem 2 was solved. This allowed the user to generate a Web page using data in the document medium. Using these functions, when different users enter same URL as input into Web browser, different information can be searched and extracted from the database or document medium and

displayed on the Web browser. However, these functions allow one to deal only with information for individual users. In practice, these functions are insufficient in an organization.

In this study, we propose the DACS Web Service, which improves the problems above. It has the functions of using information for groups of users and all users based on two functions of existing Web Service. In addition, the scheme integrates these two functions, and generates unified and personal user interfaces for each user. The personal portal for practical use in an individual organization is thus realized.

The following are improvements obtained over existing personal portal by incorporating the DACS Web Service into practical personal portals.

- (1) Given that each user can create a Web page as a personal portal that fits his/her preference without the limitation of the program on the Web server, a user interface that is easy to use and suitable for practical use is realized
- (2) Data stored not only in a database but also in a document medium, are used easily.

Next, we examine a new element that is necessary for functional extension: access control technology for the data corresponding to position and the role of the user. Role-based Access Control (RBAC) [28], [29], [30] is available for this kind of access control. As RBAC enables access control corresponding to the role of individual user, access control for groups of users as well as all users is functionally possible. RBAC can be used to add the function required in this study.

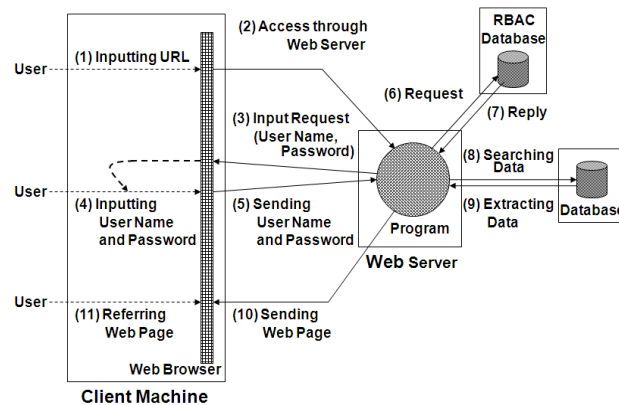


Fig. 3 Mechanism of Personal Portals with RBAC

The mechanism of personal portals when applying RBAC is shown in Fig. 3. The essential difference between Fig. 3 and Figs. 1, and 2 is that processes (6) and (7) are added in Fig. 3. These processes, which involve requesting and extracting access control information, must be performed between the gateway program on the Web server and the RBAC server that stores the access control information. When RBAC is applied to the two types of Web Service functions based on the DACS Scheme, the system configuration does not have the gateway program. Therefore, when an inquiry to a database is sent, the DACS Client must be extended to request and extract access control information from RBAC before the inquiry. In addition, because the DACS Scheme is the method used to maintain access control information in the DACS Server, but not in RBAC Server, access control information must be managed in two servers. This extension goes against the concept of the DACS Scheme, which is to control an entire network using a simple mechanism. Therefore, it is inappropriate to apply RBAC to two types of Web

Service function based on the DACS Scheme. The extension must be carried out in a form suitable to the DACS Scheme.

In Section 3, a synopsis of the following is provided to explain the DACS Web Service: the DACS Scheme and the two types of functions based on the scheme, as well as the system configuration by these functions to realize the customized personal portal. In Section 4, the actual DACS Web Service is proposed. In particular, the features that are improved relative to the two types of the conventional Web Service functions are explained. In Section 5, experimental results to confirm the possibility and an example of additional customization are described.

3. SYNOPSIS OF EXISTING DACS SCHEME

Papers in this format must not exceed twenty (20) pages in length. Papers should be submitted to the secretary AIRCC. Papers for initial consideration may be submitted in either .doc or .pdf format. Final, camera-ready versions should take into account referees' suggested amendments.

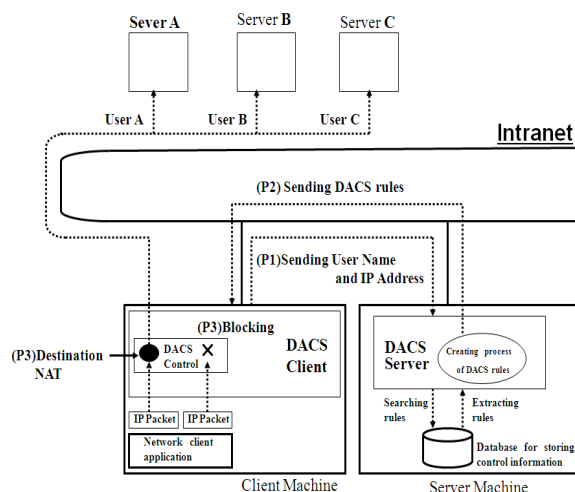


Fig. 4 Function of the DACS Scheme (1)

3.1. Functions of the DACS Scheme

First, we provide a summary of the DACS Scheme. Figs. 4 and 5 show the functions of the network services according to the DACS Scheme. The DACS Server, which is located on a server machine, refers to the server software of the DACS Scheme. The DACS Client, which is located on each client machine, is the client software of the DACS Scheme. The DACS Control and DACS SControl in the DACS Client denote the control mechanisms for communications that are sent from the client software such as a Web browser. The DACS Control controls the normal communication from client software. On the other hand, the DACS SControl converts normal communication from the client software to Virtual Private Network (VPN) [31], [32], [33], [34] communications.

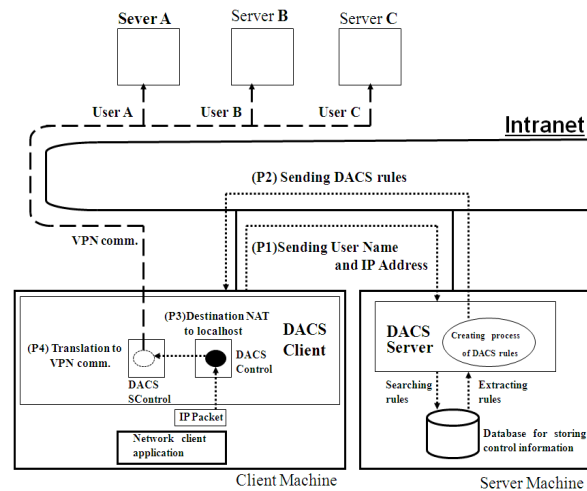


Fig. 5 Function of the DACS Scheme (2)

The DACS rules are the rules for controlling the DACS Control and DACS SControl. At the time of (a) or (b) below, the DACS rules (rules defined for each user) are distributed from the DACS Server to the DACS Client.

- (a) When the user logs into the client machine
- (b) When a notice is delivered from the system administrator

According to the distributed DACS rules, the DACS Client performs operation (1), (2) or (3) operations. Thereafter, communication control of the client machine is performed for each user who has logged in.

- (1) The destination information on the IP Packet, which is sent from the client application, is changed by Destination Network Address Translation (NAT).
- (2) The packet from the client machine, which is sent from the client application to out of the client machine, is blocked by a packet filtering mechanism.
- (3) Communication between a client machine and a network server is supported by VPN with the port forward function of Secure SHell (SSH) [35], after the destination of the communication is changed to localhost (127.0.01) by function of (1).

An example of case (1) is shown in Fig. 4. Here, communication control by the system process (P3) is performed after the system proceeds from (P1) to (P2). Thereafter, the system administrator can distribute a communication of the user who has logged in to a specified server among servers A, B, or C. An example of case (2) is also shown. When the system administrator wishes to forbid a user to use Mail User Agent (MUA) as the client software of the control target, this is achieved by blocking the IP Packet with specific destination information.

Next, an example of case (3) is shown in Fig. 5. Here, communication control by the system proceeds (P3) and then, (P4) is performed after the system proceeds from (P1) to (P2). The communication is supported by VPN, and the system administrator can distribute the VPN communication of the user who has logged in to a specific server (A, B, or C). In the DACS

Scheme, when a network service that handles personal and confidential information is controlled, the communication between a network server and client machine must be encrypted for protection. This function is used at that time.

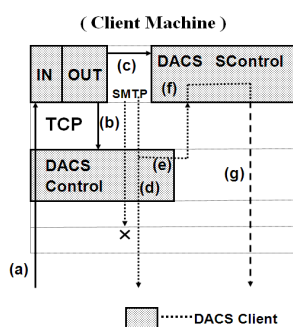


Fig. 6 Layer setting of the DACS Scheme

In order to realize the DACS Scheme, communication controls on the client machine are performed by the DACS Protocol, as shown in Fig. 6. The DACS rules are distributed from the DACS Server to the DACS Client in (a) of Fig. 6, and applied to the DACS Control and the DACS S Control in (b) and (c) of Fig. 6. Normal communication control, such as modification of the destination information or communication blocking, is performed at the network layer in (d) of Fig. 6. In case (1), as the function of destination change, the DACS rules are only applied to the DACS Control. The DACS rules consist of both destination-a (the destination IP address-a, destination port-a) before destination change and destination-b (destination IP address-b, destination port-b) after destination change. The communication sent from the client software is sent to destination-a. Thereafter, the destination of the communication is changed to destination-b by the control of destination NAT in the DACS Control, and communications are sent to the network server with this destination. In case (2), as the function of communication blocking, the DACS rules only are applied to the DACS Control similarly. The content of the DACS rules consists of destination-c (destination IP address-c, destination port-c) as the communication-blocking target. When the communication sent from the client software is destination-c, the communication is blocked by packet filtering in the DACS Control. In case (3), as the function of VPN communication, the DACS rules are applied to both the DACS Control and the DACS SControl. The content of the DACS rules consists of the following two rules.

- (r1) Rules with both destination-d (destination IP address-d, destination port-d) before destination change and destination-e (127.0.0.1, port-e) after destination change
- (r2) Rules with destination-f (destination IP address-f, destination port-f)

The DACS rules as (r1) are applied to the DACS Control and the DACS rules as (r2) are applied to the DACS SControl. Then, when a communication is supported by VPN, it is sent from (f) to (g) via (e). The VPN communication of (g) is sent by the DACS S Control. Using the port forwarding function of SSH, VPN communication which tunnels and encrypts the communication between a network server and a client machine with the DACS Client is realized. Normally, to communicate from the client software to a network server using port forwarding of SSH, it is necessary for the local host (127.0.0.1) to be indicated on that software as the communicating server. Using this function, transparent use of a client machine as a characteristic of the DACS Scheme is not failed. That the use of a client machine is transparent means that even if the configuration of the network servers is changed, the client machine can be used continuously without changing its setups. Communication control for this function is performed with the DACS SControl by SSH port forwarding. Through the use of these two functions, VPN or non-

VPN communication for each network service can be selected for each user. In the case where non-VPN communication is selected, the communication control is performed by the DACS Control, as shown in (d) of Fig. 6. In the case where VPN communication is selected, the destination of the communication is changed by the DACS Control to the localhost. Then, the port number is changed to the number assigned for each communication. Subsequently, the communicating server is determined and the VPN communication is sent by the DACS SControl, as shown in (g) of Fig. 6. In the DACS rules applied to the DACS Control, the localhost is indicated as the destination of communication. In the DACS rules applied to the DACS SControl, the network server is indicated as the destination of communication. Thereafter, by changing the content of the DACS rules applied to the DACS Control and the DACS SControl, the control in the case of VPN communication or non-VPN communication is distinguished.

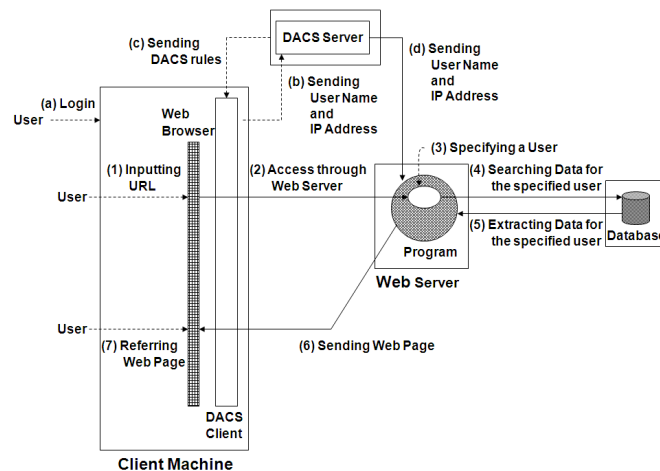


Fig. 7 Function Using Data from Database

3.2. Two Types of Web Service Functions based on DACS Scheme

In this Section, the network service corresponding to the DACS Scheme is explained. In the existing DACS Scheme, the communication control information for each user and each client machine has been maintained in the DACS rules on the DACS Server. By applying that information for communication control to the DACS Client (DACS Control) located on the client machine, communication from the client machine is controlled. As a result, the communication control mechanism on the network server is not required. However, on a network introducing the DACS Scheme, if a correspondence list of a client machine's IP address and user name is passed to the network server, it becomes possible to identify which user is sending the communication from a client machine. As a result, it becomes possible for a program on the network server to perform different processing for each user. A concrete example is the correspondence of the Web Service to the DACS Scheme. As the example, two types of Web Service functions based on the DACS Scheme are described as follows.

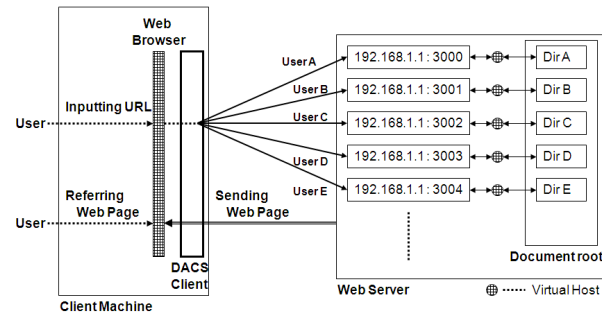


Fig. 8 Function Using Data from Document Medium

First, the function to retrieve data from the database is developed. To realize this function, the DACS Scheme must be extended, and the program on the Web server must be implemented in correspondence to the extended DACS Scheme, as shown in Fig. 7. In the existing DACS Scheme, after a user logs in to a client machine (a), the user name and IP address are sent to the DACS Server (b). Then, the DACS rules are sent back to the DACS Client (c). In the extended DACS Scheme, the user name and IP address are sent to the program on the Web server. A characteristic of the extended DACS Scheme is that the server side program on the Web server identifies the user by checking the login information and the source IP address from the client machine, and changes the processing of the program for each user. When different users access the program with the same URL, different information for each user can be searched and extracted from the database, and can be displayed on the Web browser. On the extended DACS Scheme, this new function is performed through the processing from (1) to (7).

Next, the function to retrieve data from the document medium for each user is developed. In the network with the DACS scheme, a different IP address and TCP port can be assigned for each user using the same host name. Therefore, a different document medium with the same file name on a different Web server can be referred for each user by entering the same URL into the Web browser as input. When this principle is combined with the function of a virtual host, which is equipped as a Web server, it is possible to use the Web server in the manner shown in Fig. 8. By the function of the virtual host, multiple sockets (IP address and TCP port) can be assigned to one Web server. The referred document can be changed for each socket. First, in the document root of the Web server in Fig. 8, directories (i.e., Dir A, B, C, D...) are prepared for each user. By the function of the virtual host, each directory is connected to each socket as a pair. By changing the TCP port number (3000, 3001, 3002...) for an IP address (192.168.1.1), the sockets corresponding to each directory are prepared. Next, movement on this mechanism is described. First, a user enters a URL into a Web browser as input. When the URL is entered as input by User A, the file in Dir A that is connected to the socket (192.168.1.1:3000) is referred. When User B enters the URL as input, the file in Dir B that is connected to the socket (192.168.1.1:3001) is referred. For User C, the file in Dir C that is connected to the socket (192.168.1.1:3002) is referred. When a document medium with the same name exists in each directory (i.e., Dir A, B, C...), each user views different contents by entering the same URL into the Web browser as input. From the user's point of view, the same function as that of the Web Service based on the extended DACS Scheme is realized. For the information sender, because it is possible to deliver information to the specific user by uploading the document medium to the predetermined directory, the range of information usage broadens greatly. Because the information sender can easily describe the content of the document medium, it is possible to communicate information with much expressive power and impact.

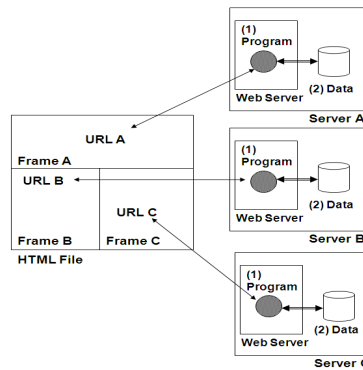


Fig. 9 Type1 System Configuration

3.3. System Configuration to Realize Personal Portals

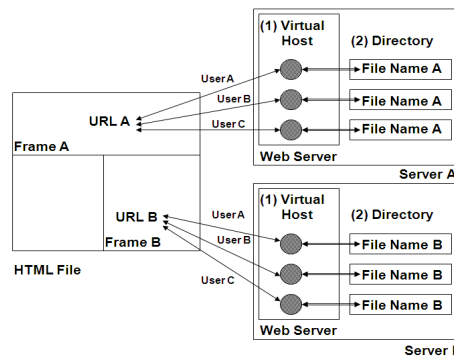


Fig. 10 Type2 System Configuration

Use of the function to retrieve data from the database enables the Type1 system configuration, which is shown in Fig. 9. First, server A-C, which have programs based on the extended DACS Scheme, as shown in (1) of Fig. 9, and data in the database as shown in (2) of Fig. 9, are distributed on the network. In the extended DACS Scheme, information related to each user is displayed on the Web browser by inputting the URL. One window of the Web browser is divided, for example, into three frames (Frames A-C). A static HTML file with each URL (URLs A-C) in each frame is created for displaying the Web page as a personal portal. The static HTML file is placed on the Web server or the client machine. When the static HTML file is opened through the Web browser, the information extracted from each server is distributed on the Web browser. In the extended DACS Scheme, the URL corresponding to each server is only incorporated in the static HTML file. Thus, when the static Web page is created, the user can easily create the customized Web page as a personal portal. As shown in Fig. 10, the Type 2 system configuration can be realized by using the function to retrieve data from the document medium. Each URL (URLs A, B) is prepared. The URL and the information that can be viewed after entering it in the Web browser as input are sent to the users. For example, the URLs for acquiring the homework in one classroom and for communicating from the office to each individual user are enumerated as different kinds of URL. By entering this URL into the Web browser as input, the file with the same name (File Name A or B) is referred. The file is stored in each directory for each user, and each user can view the stored file in the directory. Therefore, a customized personal portal with a static HTML File that is the same as Type 1 configuration can be created. As a result, by allowing both system configurations to coexist, a Web Service is realized in which a user can use information on the network regardless of the form of its storage. (Fig.11)

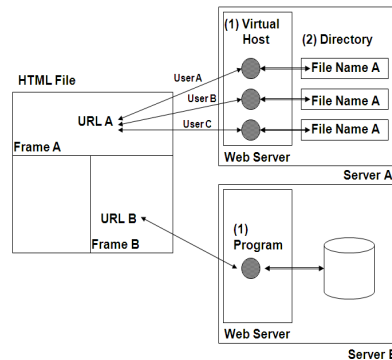


Fig. 11 Web Service by Two Functions

4. DACS WEB SERVICE

In Section III, we provided a synopsis of the DACS Scheme. In addition, we described two types of Web Service functions that could be realized in the DACS Scheme or extended DACS Scheme, and the system configuration according to these two types of functions for realizing a customized personal portal. As these two functions are used to manage the information of each user, it is insufficient to create a customized personal portal for use in one group of users. As explained in Section II, the following three functions are necessary.

- (1) A function to manage the information of each user.
- (2) A function to manage the information of user groups.
- (3) A function to manage the information of all users.

Among these three functions, the latter two do not exist. Therefore, in this section, these two types of Web Service functions are integrated after extending the function to compensate for this insufficiency. The DACS Web Service, which is achieved by this integration, is proposed to realize the personal portal. Moreover, we explain the method to acquire the data they seek.

4.1. Operational Procedures for Personal Portal

In this section, operational procedures to create a web page for a personal portal are explained. To be concrete, a system administrator's procedure and a user's procedure are described.

(System Administrator's Procedure)

- (Step1) After a system administrator assesses where data in information systems on a LAN exists, s/he make a list with contents of the data and the URLs to acquire it.
- (Step2) The list is uploaded and released on a Web Server to let users know where the data they seek exists.
- (Step3) The URL to view the lists notified to users through some methods such as the following.
 - (1) Notice on a bulletin board
 - (2) E-mail
 - (3) Notice on a Web page that many users view
 - (4) Oral notification by telephone

In Fig.12, an example of such a Web page is described. In this page, content of the data and URLs for viewing the data are described side by side on one line. When the DACS Client is installed on the client machine, the user can view this page through the function of the DACS SControl. That is, this page is accessed securely by VPN communication. When the DACS Client is not installed on the client machine, the user can not view this page. By clicking the URL of the page on a client machine with DACS Client, the user can easily acquire the data they seek. By user's procedure as follows, each user can create the web page for the personal portal.

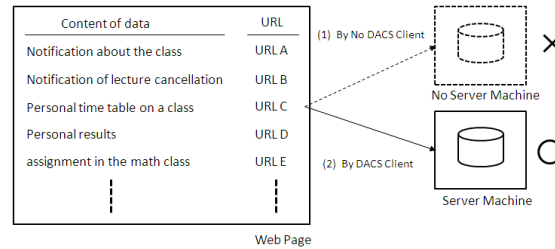


Fig. 12 Web Page to show Where Data and URLs Exist

(User's Procedure)

- (Setp4) Each user inputs access the above list by inputting the specified URL with the DACS Client installed in a client.
- (Step5) Each user views the list, and acquires the URL for getting the necessary information.
- (Step6) Each user creates a static HTML file for the personal portal by implementing the above URL.
- (Step7) Each user uploads the above HTML file.

After these steps, it is possible to view the web page for the personal portal.

4.2. Synopsis of DACS Web Service

In Fig.13, an overview of the DACS Web Service is shown. In this figure, the function to retrieve data from a database of an information system is shown as Function α , and the function to retrieve data from a document medium such as a simple text or PDF file is shown as Function β .

From here, a series of DACS Web Service's movement is explained. First, as URLs to use the DACS Web Service, three kinds of URLs as follows are prepared for Function α and Function β .

- (1) URL for acquiring data of each user
- (2) URL for acquiring data of each group
- (3) URL for acquiring data of all users

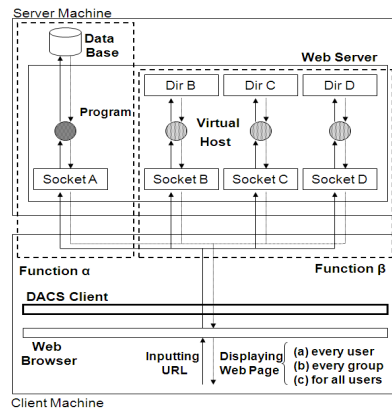


Fig. 13 DACS Web Service

By using these URLs, necessary information is acquired. The concrete functions by using these URL are described in Fig.14. By the URL (1), the existing two functions described in the Fig.7 and Fig.8 are used. By the URL (3), normal web access is done. In the case of Function α , the data is extracted from a database. For example, web access without three processes of (3),(4) and (5) in Fig.1 is given. In the case of Function β , the data is extracted from a static document medium. This is a normal web access that we do in accessing a home page on Internet. Because the above cases are realized by using the existing technology, it is not necessary to explain it especially.

URL	Function α	Function β
(1)	Acquiring data by using the function of Fig.7	Acquiring data by using the function of Fig.8
(2)	Acquiring data by using the extended function of Fig.7	Acquiring data by using the extended function of Fig.8
(3)	Normal web access by extracting data form a database	Normal web access by using a static document medium

Fig. 14 Details of Function α and Function β

However, in the case of using the URL (2), the extended functions from the functions of Fig.7 and Fig.8 are used. Therefore, technological explanation needs to be described. In Fig.15 the extended function from the function of Fig.7 is described. Processes from (4) to (7) in Fig.17 are additional processes from the function of Fig.7. First, a user name, which is sent from the DACS Sever by the process (d), is sent to an authentication server such as LDAP (5). The authentication server receives it, and returns the group name that the user belongs to by extracting it form an authentication database in process (5). The program on the Web server receives the group name, and sends it to the database in process (6). The database receives the group name, and extracts data for the group from the database. By process (7), the data is sent to the program on the Web server. The program receives the data, and sends a Web page to the Web browser on the client. In Fig.16, the extended function from the function of Fig.8 is described. This extended function is realized by changing the setups of the DACS rules. In Fig. 7, each user passes through the specified sockets and accesses the specified directory. Different multiple users do not access the same directory. However, by assigning the same DACS rules to each user belonging to the same group, the function described in Fig.16 is realized. To be concrete, when multiple users form

UserA1 to UserA5 inputs same URL to the Web Browser, the same static document medium in the directory of GroupA is accessed through the specified sockets (192.168.1.1:4001). In the same way, When multiple users form UserB1 to UserB5 inputs same URL to the Web Browser, the same static document medium in the directory of GroupB is accessed through the specified sockets (192.168.1.1:4002). After these above access, each Web page for each user's group is sent to the Web browser on the client.

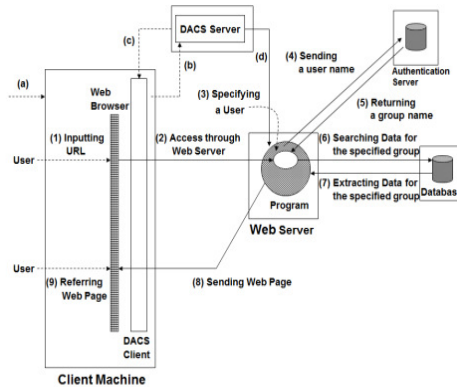


Fig. 15 Extended function of Function α

In this way, a user can use data from an information system or document medium dispersed on the network, without being aware of which function is being used. In other words, a user can use information regardless of its form or place of storage, if a user knows the URL and the kind of information acquired by that URL. Regardless of whether Function α or Function β is used, data are displayed on the Web browser after entering the URL as input. Three kinds of data which are sent to each user (a), each group of users (b), and all users, are displayed.

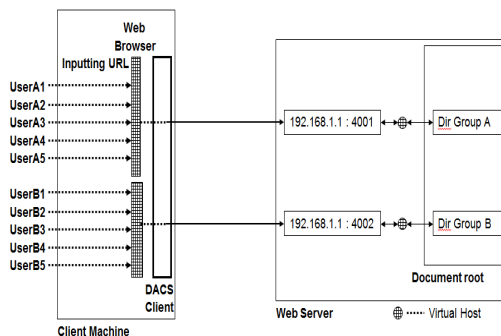


Fig. 16 Extended function of Function β

4.3. Characteristics of DACS Web Service

(1) Unifying Access Control

To manage access control information that is stored in one place, a network administrator can unify the access control for multiple Web servers distributed on the local area network. In a conventional network, access control for Web servers is performed by the mechanisms for access control that are distributed on the local area network, such as packet filtering of the fire-wall and router as well as access control of the Web server.

(2) Detailed Access Control

To change access control information that is stored in one place, a network administrator can manage the access control for the Web server according to the users and the kind of network (e.g., for office work, for students, public network). To be more specific, when a user moves multiple networks with a personal notebook computer and is permitted access to one Web server, the network administrator can refuse the user's access depending on the network configuration. In a conventional network, the same access control can be performed functionally by arranging a mechanism that combines access control by user authentication with that by use of the source IP address. However, this is a complicated and troublesome method because a network administrator needs to manage multiple Web servers individually.

(3) Change of the Referred Web Server

Using the same URL, the referred Web server can be changed according to the user and the type of network. When the Web Server, which is used by presenting the necessary information for each network (e.g., the guidance information of the network use), and the URL for information reference is decided as a promise, a user can refer to that information by automatically accessing the Web server located in that network.

5. EXPERIMENTAL RESULTS BY PROTOTYPE CONSTRUCTION

To confirm the possibility of this system, functional experiments by implementing a prototype system in Fig.17 were performed. The details of system configuration were as follows:

(1)Server Machine

CPU: Celeron M Processor340 (1.5GHz)

OS: FedoraCore3

DACS Server:

Language: Visual C++ 7.1

Database: postgresql-8.4.4-2

(2)Server Machine2

CPU: Celeron M Processor340 (1.5GHz)

OS: CentOS 6

Apache: httpd-2.2.15-5

Database: postgresql8.4.4-2

Program on Apache

Language: php-5.3.2-6

(3)Client Machine

CPU: Celeron M Processor340 (1.5GHz)

OS: WindowsXP Professional

DACS Client:

Language: VisualC++ 7.1,WinsockSPI [36],Putty [37]

(4)Others

AuthenticationServer: openldap-2.1.22-8(FedoraCore1)

DHCP Server: dhcp-4.1.1-12

DNS Server: bind-9.2.2.P3-9(FedoraCore1)

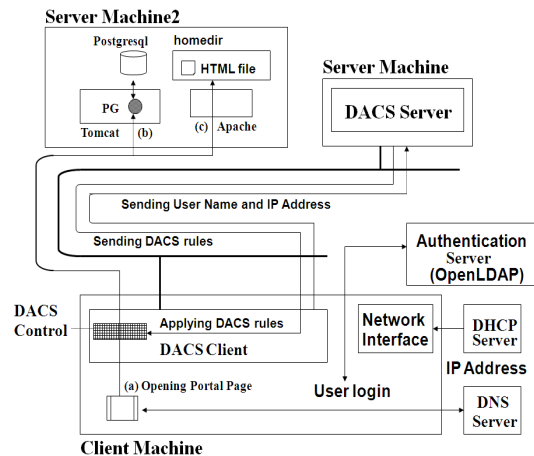


Fig. 17 Prototype

5.1. Personal Portal Creation by Using Prototype System

```

<html>
<head>
<title>Personal Portal</title>
</head>

<frameset cols="50%,50%">

<frameset rows="30%,30%,40%">
<frame src="https://www.test.ac.jp/userdataget.php" name="frame1"> ←(a)
<frame src="https://www.test.ac.jp/groupdataget.php" name="frame2"> ←(b)
<frame src="https://www.test.ac.jp/allaverageget.php" name="frame3"> ←(c)
</frameset>

<frameset rows="30%,30%,40%">
<frame src="https://www.test.ac.jp/index.html" name="frame4"> ←(d)
<frame src="https://www2.test.ac.jp/" name="frame5"> ←(e)
<frame src="https://www.test.ac.jp/allusers/index.html" name="frame6"> ←(f)
</frameset>
</frameset>
</html>
    
```

Fig.18 HTML File as Personal Portal

Using this system, the experiments were performed. First, the content of the static HTML file as a personal portal (portal page) is described in Fig.18. This personal portal is divided into 6 frames. At each point from (a) to (f), each URLs for using 6 functions described in Fig.14 were set into each frame. Next, when the portal page was opened after user A logged in on a client machine, the page shown in Fig. 19 was displayed on the Web browser. The three frames on the left were as follows. Data stored in the database for user A, i.e., personal result extracted from the table which stored the results of all students were displayed in the top frame, which was realized with the Function α used by the URL (1). Data stored in the database for group A of which user A is a member, i.e., average results of the classes that user A attended which were extracted from the table with the average results for each class were displayed in the middle frame, which was realized with the Function α used by the URL (2). Data stored in the database for all users, i.e., average results of all classes were displayed in the bottom frame, which was realized with the Function α used by the URL (3). In addition, the right three frames on the right were as follows. The content of the static HTML file stored in the home directory for user A, i.e., marked examination papers of the attendance subjects were displayed in the top frame, which was realized with the Function β used by the URL (1). The content of the static HTML file stored in the home directory which is accessible by group A of which user A is a member, i.e., teaching materials of each attendance class were displayed in the middle frame. In this frame, URLs prepared for each subject are implemented, and each user is accessible to each home directory prepared for each class in each subject. Therefore, user A became accessible to teaching materials

of each attended class. In this case, when input the link of “English”, the content of the frame was changed to the page of the attended class with links of each materials in Fig.20. These were realized with the Function β used by the URL (2). The content of the static HTML file stored in home directory for all users, i.e., notices for all students were displayed in the bottom frame, which was realized with the Function β used by the URL (3). Then, when the HTML file serving as the personal portal was opened after user B logged in on a client machine, the data related to user B was displayed on each frame of the Web browser in the same way as the above-mentioned case (Fig. 21).

<p>Result of userA</p> <table border="1"> <tr> <td>English</td> <td>Mathmatics</td> <td>Chemistry</td> <td>Physic</td> <td>Geography</td> <td>World history</td> </tr> <tr> <td>80</td> <td>85</td> <td>85</td> <td>90</td> <td>70</td> <td>75</td> </tr> </table>	English	Mathmatics	Chemistry	Physic	Geography	World history	80	85	85	90	70	75	<p>Examination paper</p> <p>English Mathmatics Chemistry Physic Geography World history</p>								
English	Mathmatics	Chemistry	Physic	Geography	World history																
80	85	85	90	70	75																
<p>Average of groupA</p> <table border="1"> <tr> <td>English</td> <td>Mathmatics</td> <td>Chemistry</td> <td>Physic</td> <td>Geography</td> <td>World history</td> </tr> <tr> <td>65</td> <td>70</td> <td>60</td> <td>75</td> <td>55</td> <td>65</td> </tr> </table>	English	Mathmatics	Chemistry	Physic	Geography	World history	65	70	60	75	55	65	<p>Teaching materials</p> <p>English Mathmatics Chemistry Physic Geography World history</p>								
English	Mathmatics	Chemistry	Physic	Geography	World history																
65	70	60	75	55	65																
<p>Average of all subjects</p> <table border="1"> <tr> <td>English</td> <td>Mathmatics</td> <td>Chemistry</td> <td>Physic</td> <td>Geography</td> </tr> <tr> <td>60</td> <td>70</td> <td>60</td> <td>65</td> <td>60</td> </tr> <tr> <td>World history</td> <td>Japanese history</td> <td>philosophy</td> <td>art</td> <td>music</td> </tr> <tr> <td>60</td> <td>70</td> <td>75</td> <td>55</td> <td>65</td> </tr> </table>	English	Mathmatics	Chemistry	Physic	Geography	60	70	60	65	60	World history	Japanese history	philosophy	art	music	60	70	75	55	65	<p>Notice (All students)</p> <p>Method of taking attendance</p> <p>About examinations</p> <p>Manner for classroamm usage</p>
English	Mathmatics	Chemistry	Physic	Geography																	
60	70	60	65	60																	
World history	Japanese history	philosophy	art	music																	
60	70	75	55	65																	

Fig.19 Portal Page for User A

In this experiment, the Function α to retrieve data from the database for each user, each group of users, and all users and the Function β to retrieve data from the document medium for these three categories, were located on the same server *machine*. However, even if each is also located on different server machines, the process is the same. As a result, it was confirmed that a user could easily create a Web page as a personal portal by implementing URLs into a static HTML file.

<p>English</p> <p>[1] [2] [3] [4] [5] [6] [7] [8] [9] [10] [11] [12] [13] [14] [15]</p>
--

Fig.20 Changed Frame

5.2. Evaluation

Based on the former section, theoretical evaluations were performed. The top frame on the left in Fig.19 and Fig.21 was realized by the Function α used by the URL (1), and the top frame on the right in Fig.19 and Fig.21 was realized by the Function β used by the URL (1). That is, the personal data for each user was displayed on both frames. In the existing personal portal described in section II, the personal data display is performed based on user identification by the user authentication mechanism and data extraction from a database based on the identification user. However, in the above implementation system, the personal data is displayed without the user authentication mechanism, and the data extraction is performed not only from a database but also form a document medium.

Result of userB <table border="1"> <thead> <tr> <th>English</th> <th>Mathematics</th> <th>Chemistry</th> <th>Philosophy</th> <th>art</th> <th>music</th> </tr> </thead> <tbody> <tr> <td>75</td> <td>60</td> <td>80</td> <td>70</td> <td>90</td> <td>95</td> </tr> </tbody> </table>	English	Mathematics	Chemistry	Philosophy	art	music	75	60	80	70	90	95	Examination paper English Mathematics Chemistry Philosophy art music								
English	Mathematics	Chemistry	Philosophy	art	music																
75	60	80	70	90	95																
Average of groupB <table border="1"> <thead> <tr> <th>English</th> <th>Mathematics</th> <th>Chemistry</th> <th>Philosophy</th> <th>art</th> <th>music</th> </tr> </thead> <tbody> <tr> <td>70</td> <td>75</td> <td>85</td> <td>60</td> <td>75</td> <td>80</td> </tr> </tbody> </table>	English	Mathematics	Chemistry	Philosophy	art	music	70	75	85	60	75	80	Teaching materials English Mathematics Chemistry Philosophy art music								
English	Mathematics	Chemistry	Philosophy	art	music																
70	75	85	60	75	80																
Average of all subjects <table border="1"> <thead> <tr> <th>English</th> <th>Mathematics</th> <th>Chemistry</th> <th>Physics</th> <th>Geography</th> </tr> </thead> <tbody> <tr> <td>60</td> <td>70</td> <td>60</td> <td>65</td> <td>80</td> </tr> </tbody> </table> <table border="1"> <thead> <tr> <th>World history</th> <th>Japanese history</th> <th>philosophy</th> <th>art</th> <th>music</th> </tr> </thead> <tbody> <tr> <td>60</td> <td>70</td> <td>75</td> <td>55</td> <td>65</td> </tr> </tbody> </table>	English	Mathematics	Chemistry	Physics	Geography	60	70	60	65	80	World history	Japanese history	philosophy	art	music	60	70	75	55	65	Notice (All students) Method of taking attendance About examinations Manner for classroom usage
English	Mathematics	Chemistry	Physics	Geography																	
60	70	60	65	80																	
World history	Japanese history	philosophy	art	music																	
60	70	75	55	65																	

Fig. 21 Portal Page for User B

Next, the middle frame on the left in Fig.19 and Fig.21 was realized by the Function α used by the URL (2), and the middle frame on the right in Fig.19 and Fig.21 was realized by the Function β used by the URL (2). That is, the data for each group was displayed on both frames. In the existing personal portal described in section II, the user identification is performed by the user authentication mechanism, and the group identification is performed based on the identification user. After that, the group data display is performed based on the identification group and data extraction from a database. But, in the implementation system, the group data is also displayed without the user authentication mechanism, and the data extraction is performed not only from a database but also from a document medium. Then, the bottom frame on the left in Fig.19 and Fig.21 was realized by the Function α used by the URL (3), and the bottom frame on the right in Fig.19 and Fig.21 was realized by the Function β used by the URL (3). That is, the data for all users was displayed on both frames. In the existing personal portal described in section II, the data for all users is displayed from a database regardless of performing the user authentication. In the implementation system, the data for all users is displayed not only from a database but also from a document medium. As common elements for these explanations, the user authentication mechanism became needless, and it became possible for data extraction not only from a database but also from a document medium. By the needless of user authentication, it became possible for each user to create the portal page in the form of implementing the URLs to it, that is, flexibility for portal page creation became higher. In addition, the preservation medium of data to use in the personal portal became included in not only a database but also a document medium.

5. CONCLUSION

In this paper, the DACS Web Service was proposed to realize practical personal portals in an organization. Existing personal portals shared the basic mechanism of generating a Web page as the personal portal using a program on a Web Server to extract data from some database in the network. However, there were two problems. To be concrete, it was impossible for a user to create a Web page as a personal portal freely and easily without restrictions of the program, and to create it from data in document mediums such as PDF and text files. To solve them, two types of Web Service functions had been previously proposed. However, these functions had yet to realize the practical personal portals in an organization. Therefore, the DACS Web Service was proposed in this paper. It is realized by extending and integrating two types of Web Service functions above without a particular mechanism for access control such as RBAC. Using the DACS Web Service, it became possible for a user to create a Web page as a personal portal freely and easily, and to use data dispersed on the network in a database or in a document medium efficiently. To show its possibility, experimental results were shown after creating the Web page

of the personal portal to exchange six basic kinds of information. In the future, we will study tools to support the operation and management of the DACS Web Service.

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